

MULTI-CRITERIA APPROACH OF DATA CENTER RESOURCES RELOCATION

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Purpose

Decreasing the total power consumption of data center servers

Problem definition

There are

Computer network data center (DC) with multiple servers S , and number of servers - N_s .

P_{max_i}, P_{min_i} - The maximum and minimum power consumption

$R_{max_i}^{CPU}, R_{max_i}^{RAM}$ - available resources (the number of cores of processors, RAM)

P_i - consumption of the server, W - bandwidth of network links

N_{vm} - consumption of the number of virtual machine (VM)

matrix X - placement of VMs,

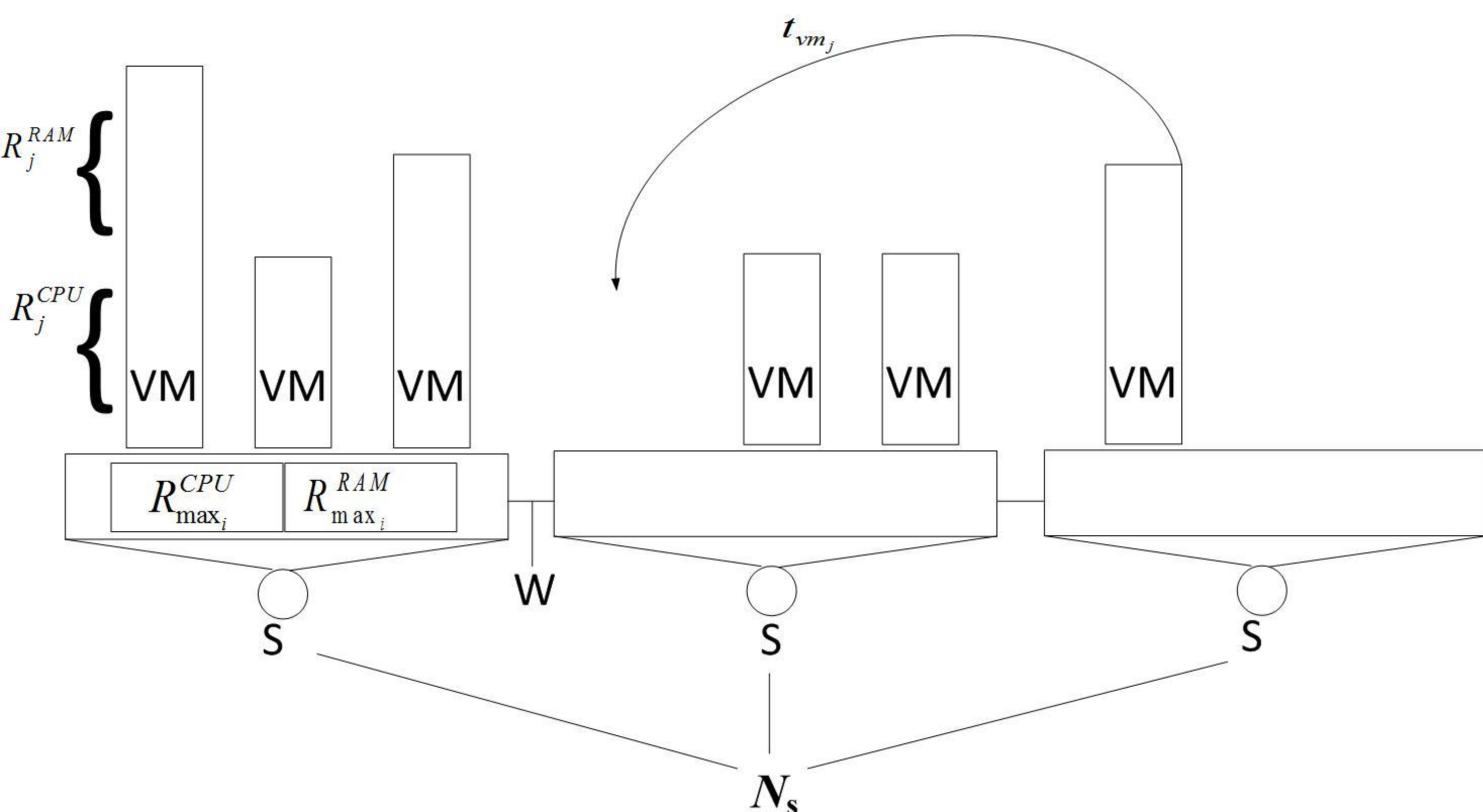
R_j^{CPU}, R_j^{RAM} - the amount of resources required for each j -th VM

Q_j^{CPU}, Q_j^{RAM} - current load of resources

$V_i(Q^{CPU}, Q^{RAM})$ - the function of the dependence of the transmitted amount of data from the current load of resources

α_i - the threshold consumption ratio of potential power, β_i - the ratio of the time of VM migration, γ - the ratio of the amount of redistributed servers

t_{vm_j} limit VM migration time



The considered problem can be formulated as below. There are multiple servers $S_i \in S$ ($i = \overline{1, N_s}$) each of which is specified by the parameters $P_{max_i}, P_{min_i}, P_i, R_{max_i}^{CPU}, R_{max_i}^{RAM}$ multiple virtual machines $VM_j \in VM$ ($j = \overline{1, N_{vm}}$) each of which is specified by the parameters $(R_j^{CPU}, R_j^{RAM}), (Q_j^{CPU}, Q_j^{RAM}), t_{vm_j}$.

Required:

matrix $X \Rightarrow$ matrix X^{final}

Method

Method for solving the problem is based on the heuristic procedure "Orderly search of options in the field of admissible solutions"

Static stage

Power consumption P_i are collected with all possible server load resource (Q_j^{CPU}, Q_j^{RAM}) Result: $P_i(Q^{CPU}, Q^{RAM})$

Dynamic stage

1. The monitoring system collects data about the reservation of each VM resources, actual resource load, SLA for each VM. The location of

each j -th VM on the i -th server is in a matrix X , where x_{ij} is a binary variable

$$x_{ij} = \begin{cases} 1, & n_{pu} VM_j \in S_i \\ 0, & n_{pu} VM_j \notin S_i \end{cases}$$

$$\text{also: } \sum_{i=1}^{N_s} \sum_{j=1}^{N_{vm}} x_{ij} = N_{vm}$$

2. For each i -server of the set S : $P_{p_i} = P_{min_i} + \sum_{j=1}^{N_{vm}} P_i(R_{vm_j}^{CPU}, R_{vm_j}^{RAM}) \cdot x_{ij}$

3. The servers will disable ($P_{p_i} = P_{min_i}$)

Then check: $P_{p_i} \leq P_{min_i} + \Delta\alpha_i$, where $\Delta\alpha_i = (P_{max_i} - P_{min_i}) \cdot \alpha_i$

Then the server is added to the set $S1$

4. For each server of set $S1$, the condition is checked: $\frac{V_i(Q_{vm_j}^{CPU}, Q_{vm_j}^{RAM})}{W} < t_{m_j} \cdot \beta_j$

Then $S2$ is formed

5. $S2$ is sorted by P_{min_i} and $\frac{P_{p_i}}{P_{min_i}}$

Then $S3$ is formed by $N_s \cdot \gamma$

The remaining servers are added to M : $M = S - S3$.

6. Equations are calculated:

$$R_{ocm_m}^{CPU} = R_{max_m}^{CPU} - \sum_{j=1}^{N_{vm}} R_{vm_j}^{CPU} \cdot x_{mj}$$

$$R_{ocm_m}^{RAM} = R_{max_m}^{RAM} - \sum_{j=1}^{N_{vm}} R_{vm_j}^{RAM} \cdot x_{mj}$$

Then, to migrate all VMs from the server k of the set $S3$ to the server m of the set M , the

following conditions must be fulfilled:

$$R_{ocm_m}^{CPU} > \sum_{j=1}^{N_{vm}} R_{vm_j}^{CPU} \cdot x_{kj}$$

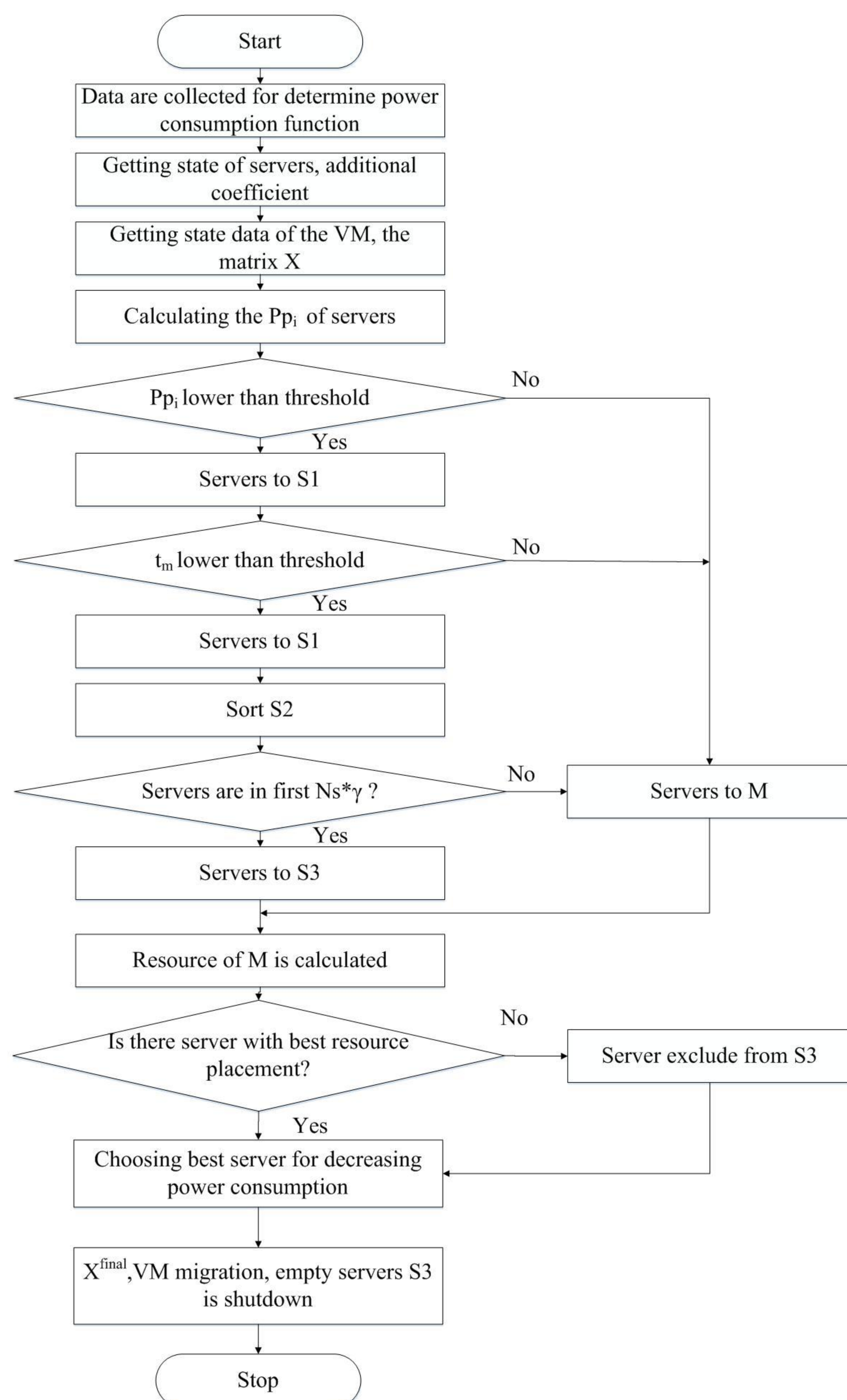
$$R_{ocm_m}^{RAM} > \sum_{j=1}^{N_{vm}} R_{vm_j}^{RAM} \cdot x_{kj}$$

$$\sum_{j=1}^{N_{vm}} P_m(R_{vm_j}^{CPU}, R_{vm_j}^{RAM}) \cdot x_{mj} < P_{min_k} + \sum_{j=1}^{N_{vm}} P_k(R_{vm_j}^{CPU}, R_{vm_j}^{RAM}) \cdot x_{kj}$$

7. The k -th server VM is migrated to the server m if the minimum value $\sum_{j=1}^{N_{vm}} P_m(R_{vm_j}^{CPU}, R_{vm_j}^{RAM}) \cdot x_{mj}$

Then server k is shutdown. Server k is excluded from $S3$. Similarly, the remaining servers of set $S3$ are checked.

The result of migrations is contained in the matrix X^{final}



Conclusions

The method allows to reduce the power consumption of the DC an amount:

$$P_{DC} - \sum_{k=1}^{N_{ds}} P_{min_k} - \sum_{k=1}^{N_{ds}} \sum_{j=1}^{N_{vm}} P_k(Q_{vm_j}^{CPU}, Q_{vm_j}^{RAM}) \cdot x_{kj} + \sum_{m=1}^{N_{ds}} \sum_{j=1}^{N_{vm}} P_m(Q_{vm_j}^{CPU}, Q_{vm_j}^{RAM}) \cdot x_{mj}^{final}$$